


Steady state and transient
state analyses of TCP and
TCP-friendly rate control
mechanism using a control
theoretic approach



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Outline



- **Introduction**
 - New transport-layer communication protocols for real-time system applications
- **Analytic model**
 - TFRC, TCP, RED
- **Steady state analysis**
 - Derive TFRC goodput, TCP goodput
- **Transient state analysis**
 - Use a control theoretic approach
- **Numerical examples**
- **Conclusion**

Background



- **Real-time applications**
 - Have been widely deployed
 - Use either UDP or TCP
- **Internet**
 - Best effort network
 - **All network applications should have a mechanism for adapting to the congestion status of a network**

UDP (User Datagram Protocol)



- Simple protocol for datagram transfer
- Doesn't have a congestion control mechanism
- **We should implement some congestion control mechanism on application layer**

TCP (Transmission Control Protocol)

- Has a congestion control mechanism
 - Adjust its packet transmission rate
- Designed for data transfer applications
 - Can tolerate a certain amount of delays
- AIMD window flow control
- Packet transmission rate fluctuates
 - **Serious problem for a real-time applications**

New transport-layer communication protocols

- **TFRC, RAP, GAIMD**
 - TCP-Friendly Rate Control
 - Rate Adaptation Protocol
 - General AIMD Congestion Control
- **Have a congestion control mechanism**
- **Realize a fairness with competing TCP flows**
- **Change the packet transmission rate smoothly**

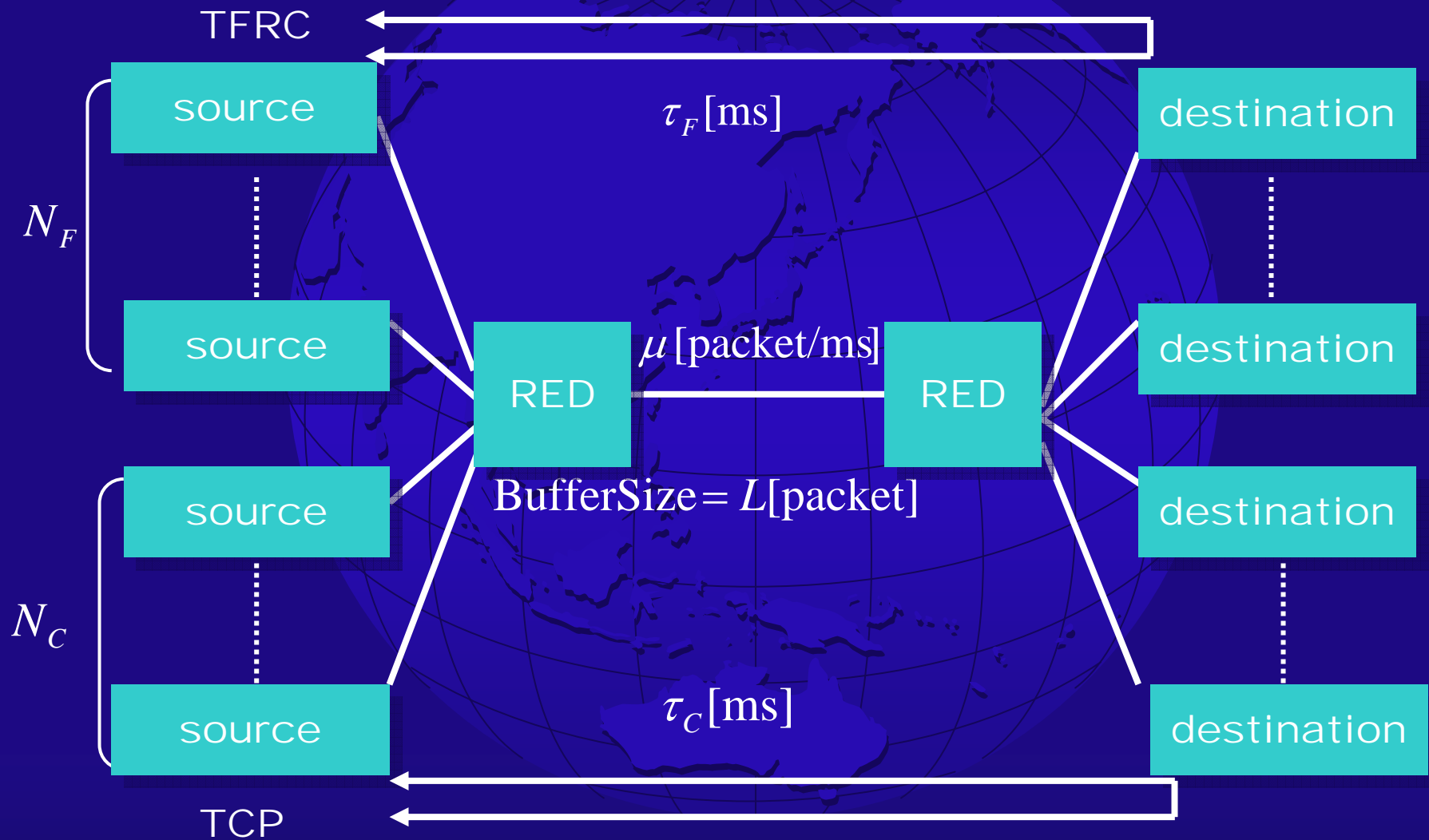
Related work

- TFRC have been studied variously
- **Steady state behavior**
 - Fairness between TFRC and TCP
 - Validity of the rate control mechanism
- **Transient state behavior**
 - Smoothness of the throughput variation
 - Responding speed to the change of the network congestion status
- **Most of these researches are based only on simulation experiments.**

Objectives

- **Model a network with TFRC and TCP connections**
 - Multiple TFRC connections
 - Multiple TCP connections
 - Single RED (Random Early Detection) router
- **Steady state analysis**
 - Derive several performance measures
 - TFRC goodput, TCP goodput, packet loss probability
- **Transient state analysis**
 - Quantitatively show convergence speed
 - Using a control theoretic approach

Analytic model



Assumption

- All TCP connections operate in their congestion avoidance phase
- Maximum window size of TCP is sufficiently larger than the bandwidth-delay product of a network
- RED routers operate appropriately
 - Average queue length of RED router is kept between min_{th} and max_{th}

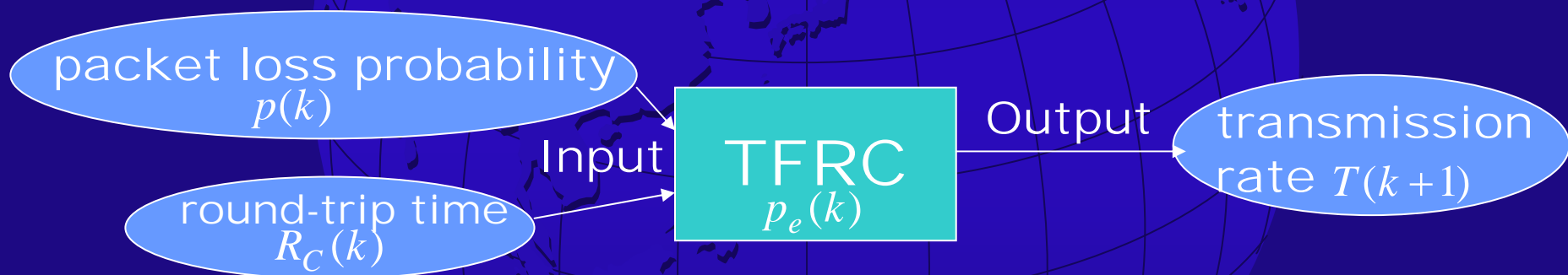
TFRC



- **Destination host**
 - measures the **loss event rate** and feeds this information back to the source host
 - loss event: one or more lost packet from a window of data
- **Source host**
 - uses feedback messages to measure the round-trip time
 - Loss event rate and round-trip time are then fed into TFRC's throughput equation
 - adjusts its transmission rate

Overview of model (1)

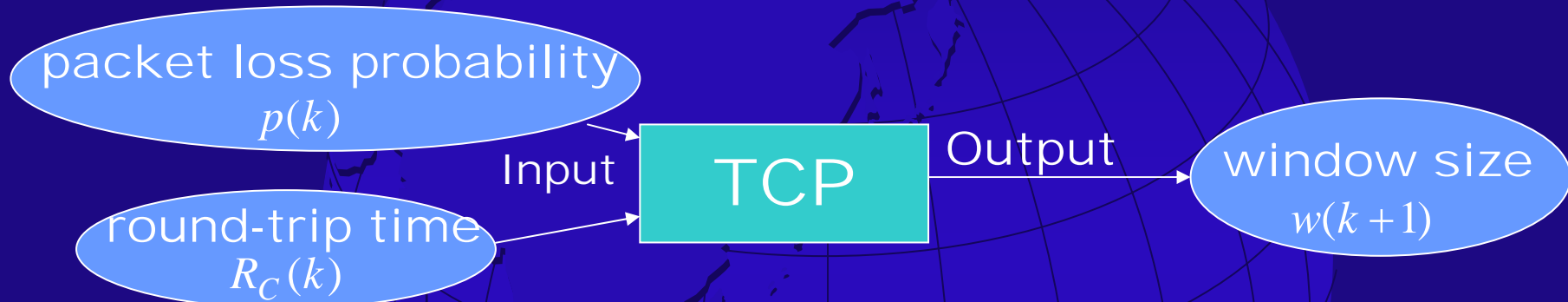
- Model TFRC, TCP, RED as discrete time systems with a time slot of Δ
- TFRC



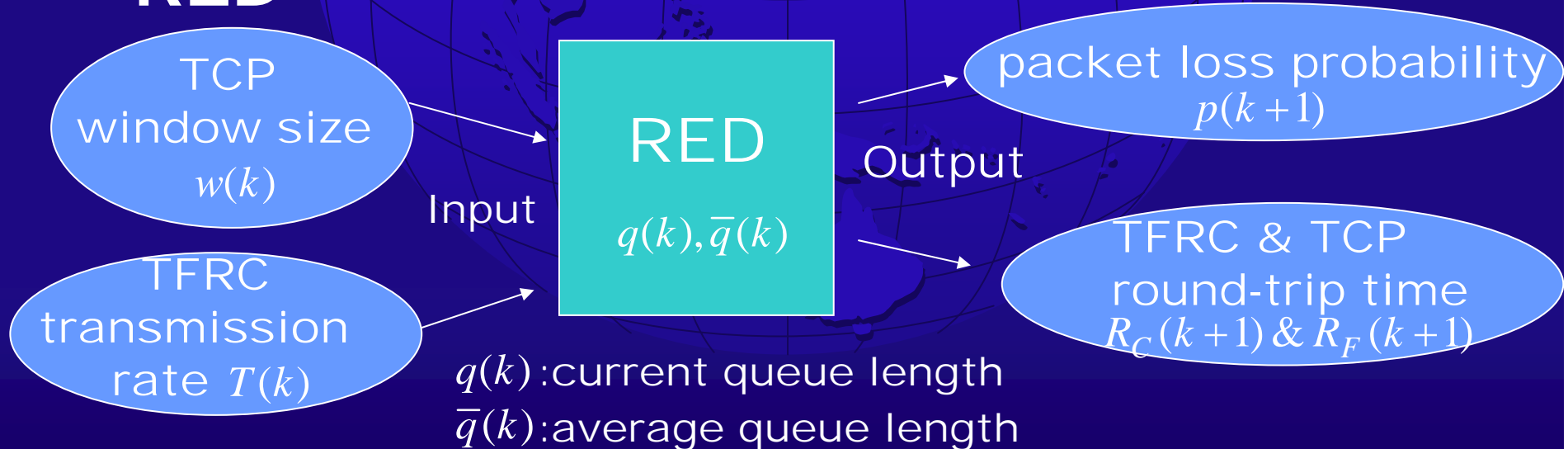
- We derive packet loss event rate $p_e(k)$, as a function of packet loss probability and round-trip time

Overview of model (2)

- **TCP**



- **RED**



Steady state analysis

- Obtain equilibrium values from models numerically
 - Equilibrium values: values in steady state
 - TCP window size w^* , round-trip time R_F^*
 - TFRC transmission rate T^* , round-trip time R_C^*
 - RED packet loss probability p^*
- Derive TFRC & TCP goodput
 - TFRC goodput: $T^* \times (1 - p^*)$
 - TCP goodput: $\frac{w^*}{R_F^*} \times (1 - p^*)$

Transient state analysis (1)

- **Define state variables & state vector**

- **State variables:** $T(k), \dots, T\left(k - \frac{R_f(k)}{\Delta}\right), w(k), \dots, w\left(k - \frac{R_c(k)}{\Delta}\right),$

- $q(k), \dots, q\left(k - \max\left(\frac{R_f(k)}{\Delta}, \frac{R_c(k)}{\Delta}\right)\right), \bar{q}(k), \dots, \bar{q}\left(k - \max\left(\frac{R_f(k)}{\Delta}, \frac{R_c(k)}{\Delta}\right)\right)$

- **State vector $x(k)$:**

- **Differences between each state variables and its equilibrium values**

$$x(k) = (T(k) - T^*, \dots, w(k) - w^*, q(k) - q^*, \dots, \bar{q}(k) - \bar{q}^*)$$

Transient state analysis (2)

- Assume TFRC notifies its source host of feedback information every M slots
- Linearize models around equilibrium points
- Obtain the transition matrix from slot k to slot $k+m$ $x(k + M) = AB^{M-1}x(k)$
 - A : state transition matrix when TFRC source receives feedback information
 - B : state transition matrix when TFRC source doesn't receive feedback information

Transient state analysis (3)

- Eigen values of AB^{M-1} determine transient state behavior
 - s : the maximum absolute eigen values of AB^{M-1} , **maximum modulus**
 - smaller s : better transient behavior
 - $s < 1$: stable
 - $s > 1$: unstable

Numerical example setting

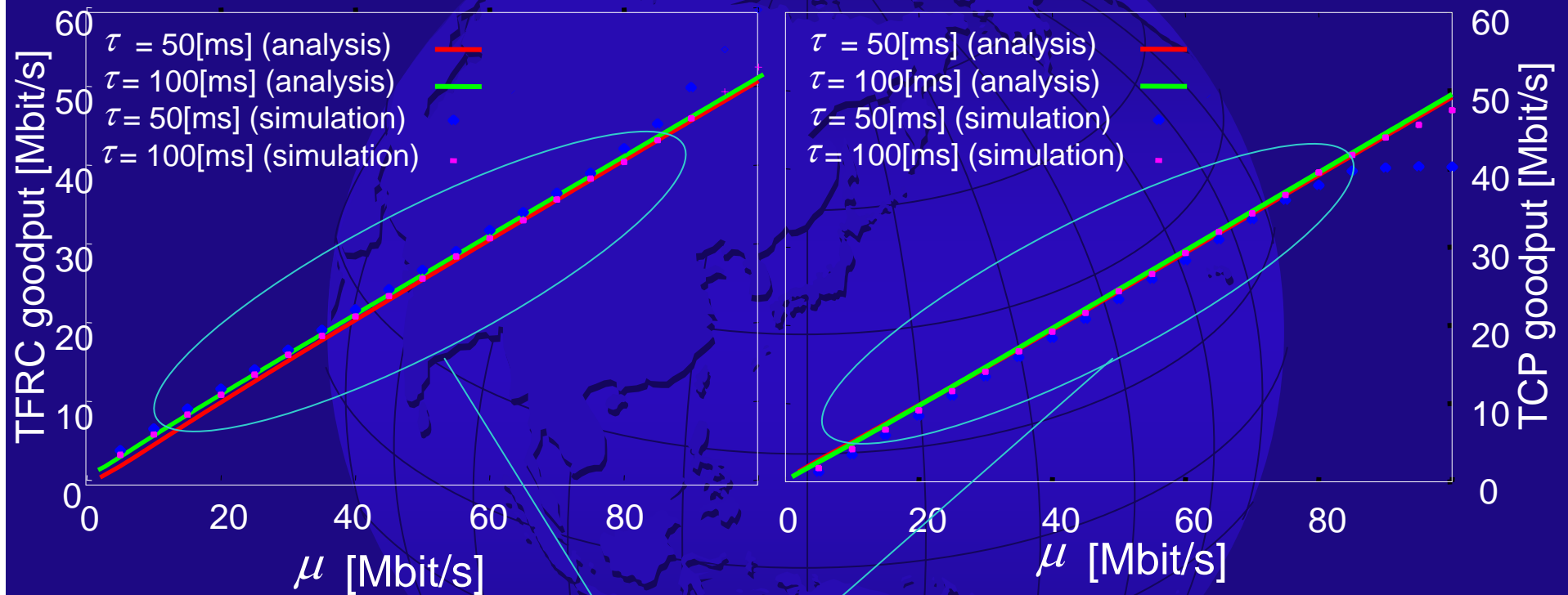
- **Analysis & simulation**

- TFRC & TCP packet size: 1000 [byte]
 - # of TFRC & TCP connections: 10, 10
 - Two-way propagation delays of TFRC & TCP are set to the equal value: $\tau_F = \tau_C = \tau$ [ms]
 - RED parameters
 - $min_{th} = 0.25\mu\tau$
 - $max_{th} = 1.25\mu\tau$
 - $L = 2.5\mu\tau$
 - $w_q = 0.002$
- $max_p = 0.1$
 L : RED buffer size [packet]
 μ : bottleneck link capacity [packet/ms]

- **Simulation**

- Simulator: ns-2
- Simulation time: 300 [s]
- # of simulation: 10

Numerical example



Good agreement

Conclusion

- **Analyze the steady state behavior of TFRC & TCP where TFRC & TCP coexist**
- **Model TFRC, TCP, RED as discrete time systems**
- **Derive TFRC & TCP goodput in steady state**
 - **Our analytic results show good agreement with simulation ones**
- **Analyze the transient state behavior of TFRC & TCP where TFRC & TCP coexist**

Thank you!

