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## Modeling of P2P File Sharing with a Level-Dependent QBD Process

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## Outline of the Talk

- Introduction and motivation
- P2P file sharing networks
  - The role of bandwidth sharing
  - File segmentation
- Analysis with level-dependent QBD
  - Scenario 1: all segments are lost
  - Scenario 2: any segment is lost (*catastrophe*)
- Numerical evaluation
- Conclusion and Outlook



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## Introduction

- Internet traffic volume has rapidly increased due to *peer-to-peer* (P2P)
- Common applications used for file sharing: eDonkey, BitTorrent, Kazaa, Winny
- Each peer acts simultaneously as *client and server*
- P2P networks form logical overlay topology on application layer



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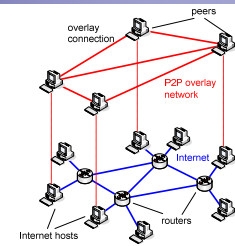
## Peer-to-Peer File-Sharing Networks

### Advantages:

- Ease of connection setup
- Load is distributed among all sharing peers  
→ *better scalability*
- Flash crowds are easily compensated

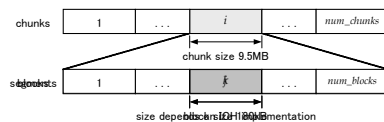
### Drawbacks:

- Data may be lost due to churn
- Content may be manipulated by peers (*pollution/poisoning*)



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## P2P System Model



- Model is similar to *eDonkey* file sharing network
- File structure consists of *chunks* and *blocks*
- Peers share blocks, corrupt chunks are discarded
- Improvement by *Intelligent Corruption Handling (segments)*



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## Download Bandwidth Sharing

- File size is  $F$  split in  $N_s = 2$  segments  
→ 3 phase system

- Asymmetrical bandwidths (ADSL)
  - Upload bandwidth  $r_u = 128$  kbps
  - Download bandwidth  $r_d = 768$  kbps

- Fair share mechanism for upload rates lead to transition rates  $\mu_1$  and  $\mu_2$ :

$$\mu_1(S, D) = \frac{1}{Z} \min \left\{ \frac{S}{D} r_u, r_d \right\} \quad \mu_2(S, D) = \frac{1}{F - Z} \min \left\{ \frac{S}{D} r_u, r_d \right\}$$

$S$ : all peers sharing segment 1  
 $D$ : all peers downloading segment 1

Download bandwidth of peer limits transmission

Upload bandwidth of peer limits transmission



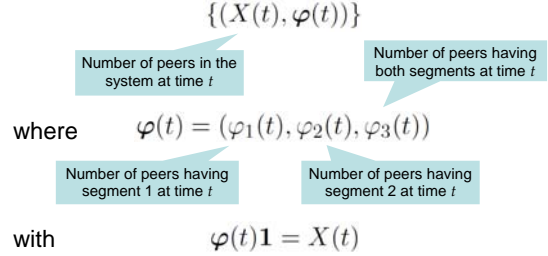
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## Analytical P2P Model

- We evaluate the *extinction probability*, when diffusion ends due to unavailability of segments
- Main contribution:**  
Proposal of algorithmically tractable analysis of level-dependent QBD process with application to P2P file sharing
- In the paper we consider two scenarios:
  - File diffusion stops when all segments are lost
  - File diffusion stops when any segment is lost

## Markovian Process Description

The system is completely described by the following Markovian process:



## Level-Dependent QBD with Catastrophes

- Extinction occurs when any segment is lost  
→ *File cannot be entirely retrieved anymore*
- Two views of the transition graph:
  - Level detailed view
  - State detailed view
- Level  $k \geq 1$  corresponds to set of states
 
$$L(k) = \{(i, j, l), i, j \in \mathbb{N}, l \in \mathbb{N}_0 \mid i + j + l = k\}$$

$$\cup \{(i, j, 0), i, j \in \mathbb{N}_0 \mid i + j = k\}$$
 and level 0 is
 
$$L(0) = \{(0, 0, 0), (n, 0, 0), (0, n, 0); n \in \mathbb{N}_0\}$$

## Level Transitions

The system may move from level  $L(k)$  to

- Level  $L(k+1)$  with rates  $A_0^{(k)}$**   
peer enters system after infinitesimal time period
- Level  $L(k)$  with rates  $A_1^{(k)}$**   
peer changes status and gets missing segment
- Level  $L(k-1)$  with rates  $A_2^{(k)}$**   
peer leaves system after infinitesimal time period
- Level  $L(0)$  with rates  $A_3^{(k)}$**   
essential peer leaves system → extinction  
e.g.  $(k-1, 1, 0)$ ,  $(1, k-1, 0)$ ,  $(0, k-1, 1)$ , or  $(k-1, 0, 1)$

## Generator Matrix

This yields the following generator matrix  $Q$  for the Markovian process  $\{(X(t), \varphi(t))\}$

$$Q = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & \dots \\ A_2^{(1)} & A_1^{(1)} & A_0^{(1)} & 0 & 0 & 0 & \dots \\ A_3^{(2)} & A_2^{(2)} & A_1^{(2)} & A_0^{(2)} & 0 & 0 & \dots \\ A_3^{(3)} & 0 & A_2^{(3)} & A_1^{(3)} & A_0^{(3)} & 0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

## State Transitions

Matrix	Transitions	Rates
$A_0^{(k)}$	$(S_1, S_2, S_3) \rightarrow (S_1+1, S_2, S_3)$	$S_1 \mu_1(L, S_2+1) + S_3 \mu_3(L, S_1+S_2+1)$
	$(S_1, S_2, S_3) \rightarrow (S_1, S_2+1, S_3)$	$S_2 \mu_2(L, S_1+1) + S_3 \mu_3(L, S_1+S_2+1)$
$A_1^{(k)}$	$(S_1, S_2, S_3) \rightarrow (S_1-1, S_2, S_3+1)$	$\mu_1(S_2+S_3, S_1)$
	$(S_1, S_2, S_3) \rightarrow (S_1, S_2-1, S_3+1)$	$\mu_2(S_1+S_2, S_2)$
$A_2^{(k)}$	$S_1 > 1 \vee S_3 > 0: (S_1, S_2, S_3) \rightarrow (S_1-1, S_2, S_3)$	$S_1 d$
	$S_2 > 1 \vee S_3 > 0: (S_1, S_2, S_3) \rightarrow (S_1, S_2-1, S_3)$	$S_2 d$
	$(S_1 > 1 \wedge S_2 > 1) \vee S_3 > 0: (S_1, S_2, S_3) \rightarrow (S_1, S_2, S_3-1)$	$S_3 d$
$A_3^{(k)}$	$S_3 > 0: (0, S_2, 1) \rightarrow (0, k, 0)$	$d$
	$S_2 > 0: (1, S_2, 0) \rightarrow (0, k, 0)$	$d$
	$S_1 > 0: (S_1, 0, 1) \rightarrow (k, 0, 0)$	$d$
	$S_1 > 0: (S_1, 1, 0) \rightarrow (k, 0, 0)$	$d$
	$(0, 0, 1) \rightarrow (0, 0, 0)$	$d$

## Extinction Probability

- The extinction probability is computed starting from state  $(0, 0, 1)$ , i.e., a single sharing peer
- We define two types of matrices:
  - $(C_0^{(k)})_{ij}$  probability to reach level  $L(0)$  in phase  $j$ , given that the process starts at  $L(k)$  in phase  $i$  and avoids  $L(1), L(2), \dots, L(k-1)$
  - $(G_k)_{ij}$  probability to reach level  $L(k-1)$  in phase  $j$  given that the process starts from level  $L(k)$  in phase  $i$
- The quantity of interest is thus

$$C_0^{(1)} = G_1$$

## Extinction Probability (2)

We define the following notations:

$$Q_0^{(k)} = (-A_1^{(k)})^{-1} A_0^{(k)} \quad \text{for the probability that starting from } L(k), \text{ the next level-transition observed is to level } L(k+1)$$

$$Q_2^{(k)} = (-A_1^{(k)})^{-1} A_2^{(k)} \quad \text{for the probability that starting from } L(k), \text{ the next level-transition observed is to level } L(k-1)$$

$$Q_3^{(k)} = (-A_1^{(k)})^{-1} A_3^{(k)} \quad \text{for the probability that starting from } L(k), \text{ the next level-transition observed is to level } L(0)$$

## Extinction Probability (3)

- We now study matrices  $G_k$  for  $k \geq 2$  satisfying:

$$G_k = \underbrace{Q_2^{(k)}}_{L(k) \rightarrow L(k-1)} + \underbrace{Q_0^{(k)}}_{L(k) \rightarrow L(k+1)} \underbrace{G_{k+1}}_{L(k+1) \rightarrow L(k)} \underbrace{G_k}_{L(k) \rightarrow L(k-1)}$$

one level transition
one or more level transitions

- This may be rewritten as

$$G_k = [I - Q_0^{(k)} G_{k+1}]^{-1} Q_2^{(k)}, \quad k \geq 2$$

- The **Logarithmic Reduction** algorithm allows computation of any  $G_k, k \geq 2$ , but it is sufficient to compute  $G_M (M \geq 2)$
- Different for  $k = 1$ :

$$C_0^{(1)} = G_1 = \underbrace{Q_2^{(1)}}_{L(1) \rightarrow L(0)} + \underbrace{Q_0^{(1)}}_{L(1) \rightarrow L(2)} \left[ \underbrace{G_2}_{L(2) \rightarrow L(1)} \underbrace{G_1}_{L(1) \rightarrow L(0)} + \underbrace{C_0^{(2)}}_{L(2) \rightarrow L(0), \setminus L(1)} \right]$$

$L(0)$  is reached without returning to  $L(1)$

## Extinction Probability (4)

More generally, to know  $C_0^{(k)}$  we first need to know  $C_0^{(k+1)}$  via the general recursion

$$C_0^{(k)} = \underbrace{Q_2^{(k)}}_{L(k) \rightarrow L(0)} + \underbrace{Q_0^{(k)}}_{L(k) \rightarrow L(k+1)} \left[ \underbrace{G_{k+1}}_{L(k+1) \rightarrow L(k)} \underbrace{C_0^{(k)}}_{L(k) \rightarrow L(0), \setminus L(1), \dots, L(k-1)} + \underbrace{C_0^{(k+1)}}_{L(k+1) \rightarrow L(0), \setminus L(1), \dots, L(k)} \right]$$

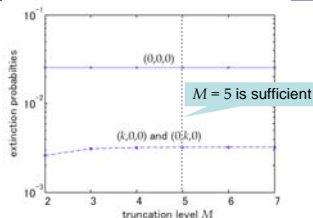
which may be rewritten as

$$C_0^{(k)} = [I - Q_0^{(k)} C_{k+1}]^{-1} [Q_2^{(k)} + Q_0^{(k)} C_0^{(k+1)}]$$

→ Truncation of the QBD after level  $L(M)$ , so that

$$C_0^{(M)} = Q_3^{(M)}$$

## Accuracy of Truncation Level



- File size  $F = 9.28$  MB, constant death rate  $d = 10^{-2}$
- Whole file extinction  $(0, 0, 0)$  unaffected by  $M$
- Truncation at  $M = 5$  provides sufficient accuracy

## Extinction Probability (5)

We finally obtain the following system which provides us the absorption probability  $C_0^{(1)} = G_1$

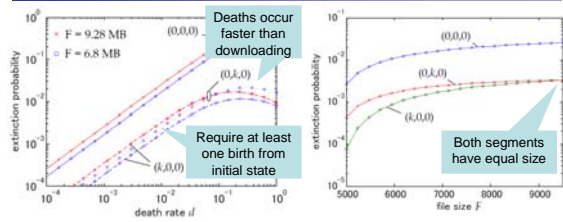
$$C_0^{(M)} = Q_3^{(M)},$$

$$C_0^{(M-1)} = [I - Q_0^{(M-1)} G_M]^{-1} [Q_3^{(M-1)} + Q_0^{(M-1)} C_0^{(M)}]$$

⋮

$$G_1 = [I - Q_0^{(1)} G_2]^{-1} [Q_2^{(1)} + Q_0^{(1)} C_0^{(2)}]$$

## Numerical Evaluation



- Smaller size of second segment leads to lower extinction probability
- Extinction probability decreases when  $d \gg \mu_1(1,1) + \mu_2(1,1)$

## Conclusion

- We proposed algorithmically tractable analysis of level-dependent QBD process
- Application to P2P file sharing network with segmentation
  - Peers should have long sojourn time to maintain low extinction probability
  - Content provider should offer incentives
- Future work includes derivation of further (transient) performance measures and more detailed user behavior (impatience, pollution, ...)