Router buffer re-sizing for short-lived TCP flows

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Research Backgrounds: Router Buffers

- Router buffer size affects:
  - Packet loss ratio in the network when the link is congested
  - Utilization of output link bandwidth
  - Packet transmission delay caused by queuing at the buffer

Two disciplines for router buffer size (1)

- Normal discipline
  - Based on Bandwidth-Delay Product (BDP) of the network
  - Bandwidth: Output link bandwidth
  - Delay: Average Round-Trip Time (RTT) of TCP connections passing through the link
  - Typical value: 250 msec
  - Determined from window size control mechanism of TCP
    - It increases window size additively when no packet loss occurs, and halves it when packet loss occurs
    - This buffer size can avoid the underutilization of the output link

Two disciplines for router buffer size (2)

- Problem in normal discipline
  - Large cost for implementing large buffers
    - 2.5Gbits buffer is necessary for a link of 10Gbps bandwidth and 250msec delay
    - Large power consumption, board size, and monetary cost
- sqrtN discipline
  - Buffer size can be decreased to bandwidth-delay product divided by the square-root of the number of TCP connections passing through the link
  - Without underutilization of output link bandwidth
  - When more than 500 TCP connections exists in the network
  - For example, buffer size decreases to 1/100 when 10,000 connections exist

Research Objectives

- Compare two disciplines (normal and sqrtN) by simulation experiments
  - Confirm the results in [2] in terms of link utilization
  - Performance of long-lived and short-lived TCP connections passing through the link
  - Evaluations in various network environment
    - Access link bandwidth
    - bottleneck link bandwidth, bottleneck link propagation delay
    - traffic volume (number of TCP connections)

Simulation environment

- Traffic models
  - P2P: each sender TCP has infinite data to transmit
  - Web: Transmission data size and transmission interval are determined by Web traffic model

Basic performance \((D=90\text{ msec}, C=100\text{ Mbps})\)

- Normal discipline provides high link utilization and small packet loss ratio
- sqrtN discipline also provides high utilization when number of flows is large enough
  - Conformance to results in [2]
  - But, packet loss ratio is significantly large

Performance of short-lived TCP connections \((C=100\text{ Mbps})\)

- Packet loss ratio in sqrtN discipline is always larger than that in normal discipline
- When propagation delay is small, larger packet loss ratio in sqrtN discipline does not affect the transmission time
- When propagation delay is large, however, the transmission delay is significantly deteriorated

Effect of access link bandwidth \((D=90\text{ msec}, C=100\text{ Mbps})\)

- Large access link bandwidth increases packet loss ratio, especially in sqrtN discipline
  - Transmission time in sqrtN discipline is up to 50% larger
  - However, small data (<100KB) transmission is not affected by the large packet loss ratio

Effect of bottleneck link bandwidth

- In sqrtN discipline, packet loss ratio never decreases to zero, although the link is underutilized
  - Because bursty packet arrival can not be absorbed at the smaller buffer
- Link utilization of sqrtN discipline is smaller than that of normal discipline
  - Using small buffer in non-congested network decrease the link utilization

Conclusion

- We tested the performance of sqrtN discipline for router buffer sizing
  - It can maintain the link utilization when there is enough traffic volume
  - But, it degrades the link utilization in non-congested network
  - It would degrade short-lived TCP performance due to large packet loss ratio
  - It may be useful only when the transmission data size is 50~100 Kbytes or when the propagation delay between the sender and the receiver hosts is significantly small
  - Otherwise, we should use large buffers
- Future work
  - The effect of pacing TCP on the buffer sizing problem